

Los complejos cúbicos $CAT(0)$ en la robótica

(y en muchos otros lugares)

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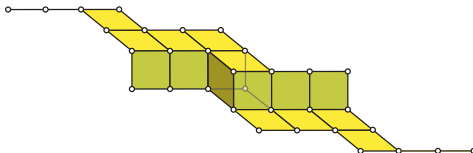
Trabajos con:

- Megan Owen (Waterloo), Seth Sullivant (NCSU)
- Rika Yatchak (SFSU/Linz), Tia Baker (SFSU)
- Diego Cifuentes (Los Andes/MIT), Steven Collazos (SFSU)



Las dos cosas que quiero decir hoy:

1. Hay muchos complejos cúbicos $CAT(0)$ “en la naturaleza”.
2. Los complejos cúbicos $CAT(0)$ tienen una estructura muy elegante y muy útil.



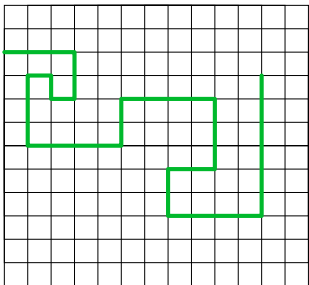
1. MOTIVATION.

Moving robots.

A robotic snake can move:

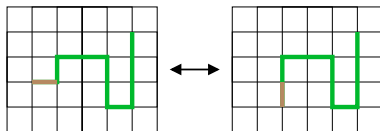
1. the head or tail or
 2. a joint
- without self-intersecting.

Snake:

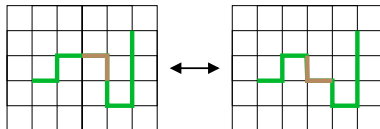


Moves:

1:



2:



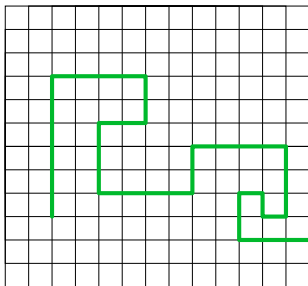
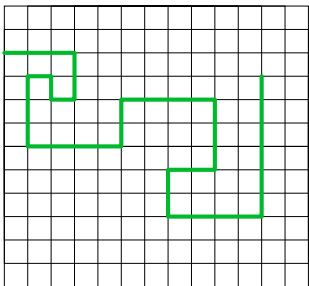
One motivation: moving robots.

How do we move this robotic snake (optimally) using these moves from one position to another one?

Position 1



Position 2



Motivación: una pregunta más fácil.

Cómo llego a la universidad?

Plaza de Nariño



Universidad de Nariño

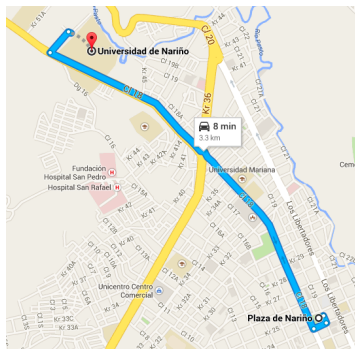


Motivación: una pregunta más fácil.

¿Cómo llego a la universidad, **de manera óptima**?

¡Con un mapa!

(Ojo: ¿Óptima en qué sentido?)



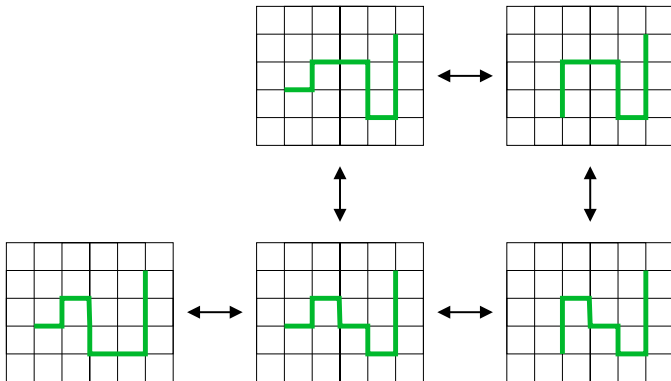
Hagamos lo mismo.

Constuyamos un mapa de las posibles posiciones del robot.

Motivation: back to moving robots.

Let's build a map of all possible positions of the robot.

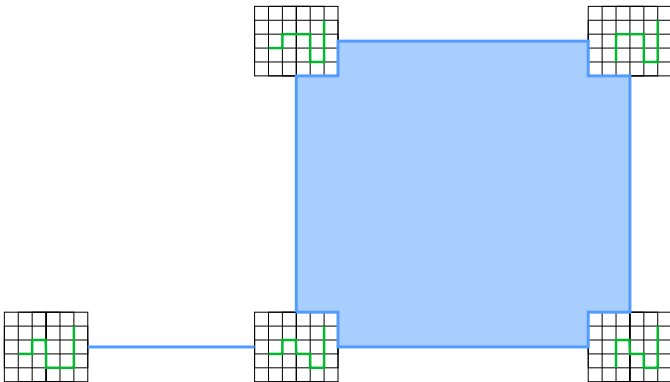
A small piece: (discrete model)



Motivation: moving robots.

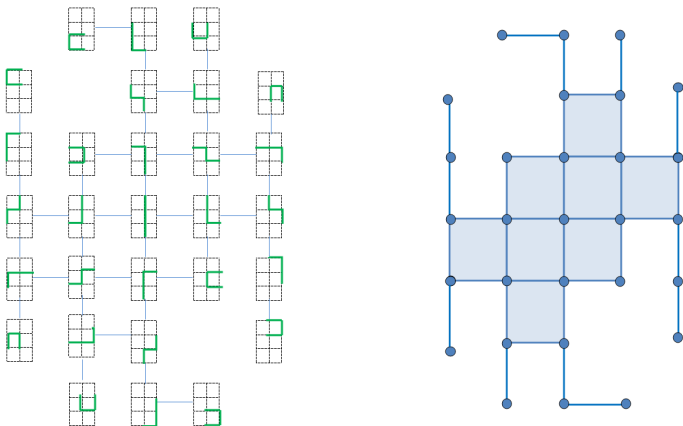
Let's build a map of all possible positions of the robot.

A small piece: (continuous model)



Motivation: moving robots.

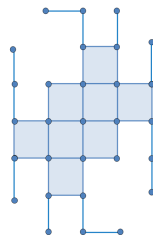
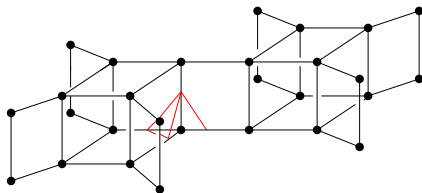
Let's build a map of all possible positions. A complete example:



A CAT(0) cube complex!

How can we understand them? Navigate them?

Motivation: moving robots.



How can we understand CAT(0) cube complexes?
How should we navigate them?

Obstacles:

- High dimension.
- Complicated ramification.
- Too many vertices.

This is what we need to overcome.

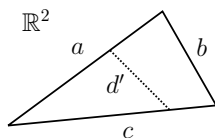
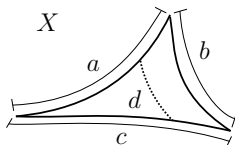
2. PRELIMINARIES. CAT(0) spaces

A metric space X is **CAT(0)** if it has non-positive curvature everywhere, in the sense that triangles in X are “thinner” than flat triangles. Roughly, it is “saddle shaped”.

More precisely, we require:

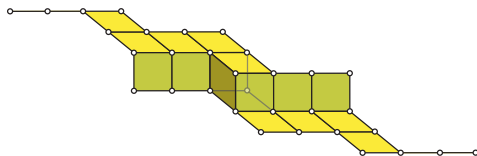
- There is a unique geodesic path between any two points of X .
- (**CAT(0) inequality**) Consider any triangle T in X and a *comparison triangle* T' of the same sidelengths in the Euclidean plane \mathbb{R}^2 . Consider any chord (of length d) in T and the corresponding chord (of length d') in T' . Then

$$d \leq d'.$$



PRELIMINARIES. Cube complexes

A **cube complex** is a space obtained by gluing cubes (of possibly different dimensions) along their faces.

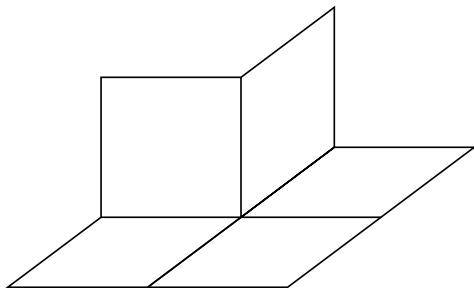


(Like a simplicial complex, but the building blocks are cubes.)

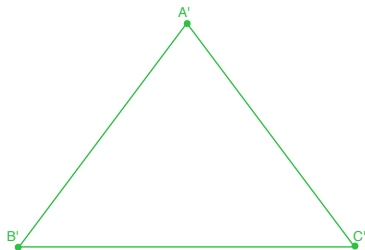
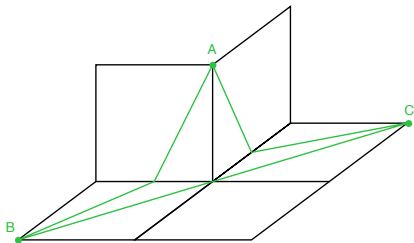
Metric: Euclidean inside each cube.

We are interested in **cube complexes which are CAT(0)**.

Example. Five squares glued around a corner.



Example. Five squares glued around a corner.

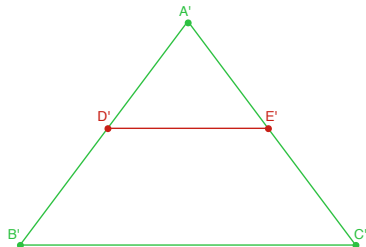
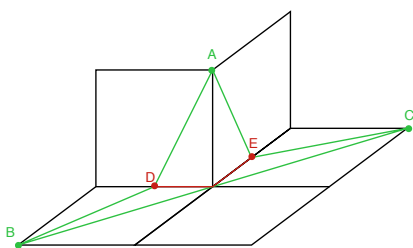


$$\begin{aligned} |AB| &= |AC| = \sqrt{5} \\ |BC| &= 2\sqrt{2} \end{aligned}$$

→

$$\begin{aligned} |A'B'| &= |A'C'| = \sqrt{5} \\ |B'C'| &= 2\sqrt{2} \end{aligned}$$

Example. Five squares glued around a corner.



$$|AB| = |AC| = \sqrt{5} \quad \rightarrow \quad |A'B'| = |A'C'| = \sqrt{5}$$

$$|BC| = 2\sqrt{2} \quad \quad \quad |B'C'| = 2\sqrt{2}$$

$$|DE| = 1 < \sqrt{2} = |D'E'|.$$

This triangle is thin. (But: I still need to test many chords.)

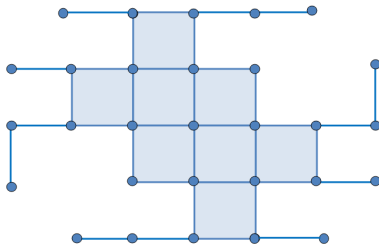
This space is CAT(0). (But: I still need to test many triangles.)

Not so practical!

3. EXAMPLES.

Example 1. Robot motion planning

State complex. vertices = positions. edges = moves.
cubes = “physically independent” moves.



Theorem (GP) This is **often** a CAT(0) cube complex.

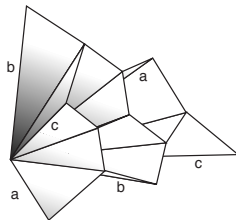
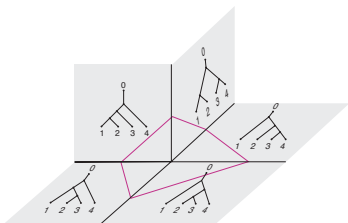
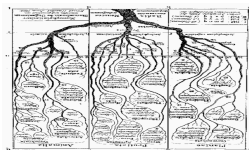
This works **very** generally for many **reconfiguration systems**, where we change vertex labels on a graph using local moves.

Example 2. Phylogenetic trees (Billera, Holmes, Vogtmann):

Goal: Predict the evolutionary tree of n current-day species/languages/....

Approach:

- Build a space T_n of all possible trees.
- Study it, navigate it.



Thm. (BHV)

T_n is a CAT(0) cube complex.

Cor. T_n has unique geodesics.

Cor. "Average" trees exist.

Example 3. Geometric Group Theory.

A **right-angled Coxeter group** is a group of the form

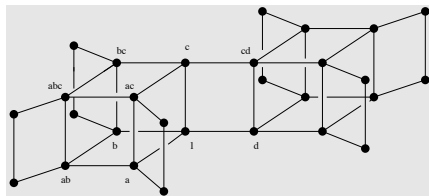
$$W(G) = \langle v \in V \mid v^2 = 1 \text{ for } v \in V, (uv)^2 = 1 \text{ for } uv \in E \rangle$$

Example: $a^2 = b^2 = c^2 = d^2 = 1$

$$(ab)^2 = (ac)^2 = (bc)^2 = (cd)^2 = 1$$



Thm. (Davis) Right-angled Coxeter groups are CAT(0):
 $W(G)$ acts “very nicely” on a CAT(0) cube complex $X(G)$.



Use the geometry of $X(G)$ to study the group $W(G)$; e.g.,

- If a group G is CAT(0), the “word problem” is easy for G .

4. CHARACTERIZATIONS.

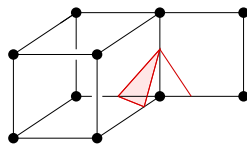
Which cube complexes are CAT(0)?

In general, CAT(0) is a subtle condition; but for cube complexes:

1. Gromov's characterization.

Theorem. (Gromov, 1987)

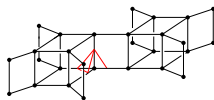
A cube complex is CAT(0) if and only if it is **simply connected** and the link of every vertex is a **flag** simplicial complex.



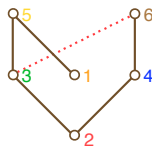
Δ flag: if the 1-skeleton of a simplex T is in Δ , then T is in Δ .

Characterizations: Which cube complexes are CAT(0)?

2. Our characterization.



Theorem. (A-Owen-Sullivan 08)
 (Pointed) CAT(0) cube complexes are in
 bijection with posets with inconsistent pairs.



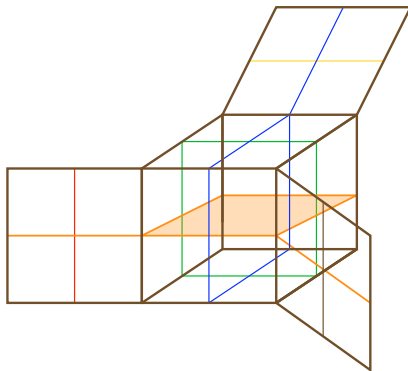
PIP: A poset P and a set of “inconsistent pairs” $\{x, y\}$, with
 x, y inconsistent, $y < z \rightarrow x, z$ inconsistent.

Theorem. (A-Owen-Sullivant 08)
 (Pointed) CAT(0) cube complexes are in
 bijection with posets with inconsistent pairs.

Sketch of proof.

Idea: CAT(0) cube complexes “look like” distributive lattices.
 So imitate Birkhoff’s bijection: distributive lattices \leftrightarrow posets

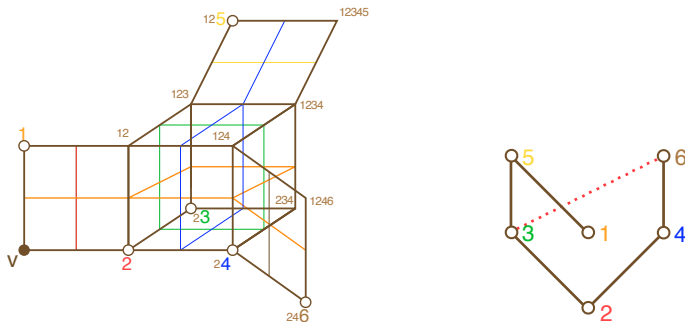
“ \rightarrow ”: X has **hyperplanes** which split cubes in half. (Sageev)



Theorem. (A. - Owen - Sullivant 08)

(Pointed) CAT(0) cube complexes are in bijection with posets with inconsistent pairs.

Bijection. “ \rightarrow ”: Fix a “home” vertex v .



If i, j are hyperplanes, declare:

$i < j$ if one needs to cross i before crossing j
 i, j inconsistent if it is impossible to cross them both.

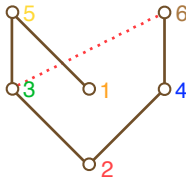
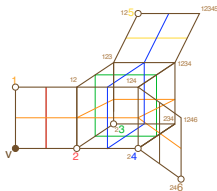
Remark.

Sageev (95) and Roller (98) obtained a different combinatorial description. Which one is more useful depends on the context.

Let's see some applications.

Application 1. Embeddability conjecture.

Conjecture. (Niblo, Sageev, Wise) Any d -dimensional interval in a CAT(0) cube complex can be embedded in the cubing \mathbb{Z}^d .



Proof. (AOS 08)

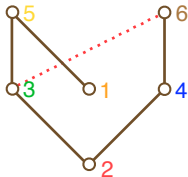
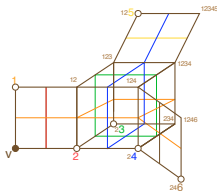
Dilworth already showed (in 1950!) how to embed $J(Q)$ in \mathbb{Z}^d :

- Write Q as a union of d disjoint chains. (Example: 246, 35, 1)
- "Straighten" the cube complex along each chain. \square

(Proof also by Brodzki, Campbell, Guentner, Niblo, Wright (08).)

Application 1. Embeddability conjecture.

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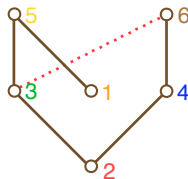
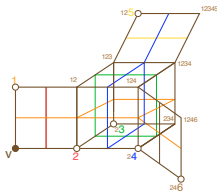
(Proof also by Brodzki, Campbell, Guentner, Niblo, Wright (08).)

Application 2. All CAT(0) cube complexes are “robotic”.

Theorem. (Ghrist-Peterson 07)

Every CAT(0) cube complex can be realized as a state complex.

Their proof is indirect.



Alternative proof. (AOS 10)

Root $X \rightarrow$ poset with inconsistent pairs P .

A “virus robot” takes over the poset P .

It can take over a new cell q if and only if:

- o it already took over all elements $p < q$, **and**
- o it hasn't taken over any elements inconsistent with q .

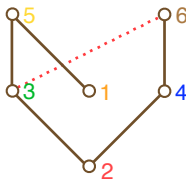
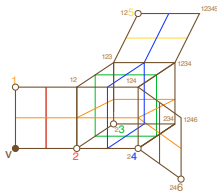
Then X is the state complex for this robot. \square

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Application 3. The Hopf algebra of CAT(0) cube complexes.

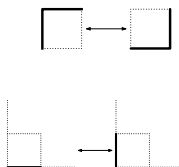
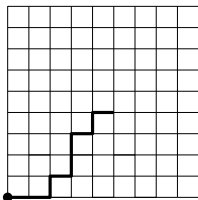
Theorem. (A. - Cifuentes - Collazos 12)

CAT(0) cube complexes have the structure of a Hopf algebra.
There is an elegant formula for the antipode.

$$\begin{aligned}
 S\left(\begin{array}{c} \bullet \\ | \\ \text{---} \text{---} \text{---} \\ | \\ \bullet \end{array} \right) &= - \begin{array}{c} \text{---} \text{---} \text{---} \\ | \\ \bullet \end{array} \\
 &+ \begin{array}{c} \text{---} \text{---} \\ | \\ \bullet \end{array} + \begin{array}{c} \text{---} \text{---} \\ | \\ \bullet \end{array} + \begin{array}{c} \text{---} \text{---} \\ | \\ \bullet \end{array} \\
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 \end{aligned}$$

The diagram shows the antipode map S applied to a diamond-shaped graph with three vertices labeled 1, 2, and 3. The result is a sum of 12 terms, each representing a different permutation of the vertices. The terms are arranged in a grid and include signs (+ or -) and labels in green boxes (e.g., 123, 12-3, 1-23, 1-2-3, 2-3-1, 23-1, 23-1, 13-2, 1-2-3, 1-3-2, 1-3-2, 2-1-3, 2-1-3, 3-1-2, 3-1-2, 3-2-1, 3-2-1). The vertices are colored: 1 is red, 2 is green, and 3 is blue.

Application 4.1. Pinned-down robotic arm in a square grid.

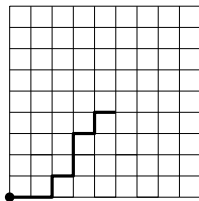


Theorem. (A.-Baker-Yatchak, 2012) The state complex is a CAT(0) cubical complex. Its PIP (“remote control”) is as shown.



Complex of 2^n states in $\frac{n}{2}$ dim. $\rightarrow \sim \frac{1}{2}n^2$ “buttons”.

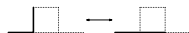
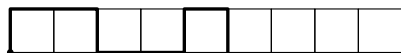
Application 4.1. Pinned-down robotic arm in a square grid.



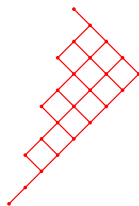
Corollary. (A.-Baker-Yatchak, 2012) Let $q_{n,d}$ be the number of d -cubes in the state complex for the robotic arm of length n . Then

$$\sum_{n,d \geq 0} q_{n,d} x^n y^d = \frac{1 + xy}{1 - 2x - x^2 y}.$$

Application 4.2. Pinned-down robotic arm in a strip.

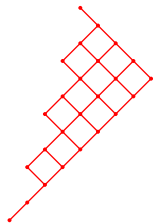


Theorem. (A.-Baker-Yatchak, 2012) The state complex is a CAT(0) cubical complex. Its PIP (“remote control”) is as shown.



Complex of $F_n \sim \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2} \right)^n$ states in $\frac{n}{3}$ -dim $\longrightarrow \sim \frac{n^2}{4}$ buttons.

Application 4.2. Pinned-down robotic arm in a strip

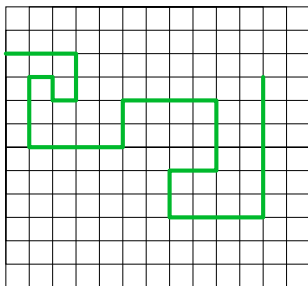


Corollary. (A.-Baker-Yatchak, 2012) Let $s_{n,d}$ be the number of d -cubes in the state complex for the robotic arm of length n .

Then

$$\sum_{n,d \geq 0} s_{n,d} x^n y^d = \frac{1 + x + xy + x^2 y}{1 - x - x^2 - x^3 y}.$$

Application 4.3. Non-pinned-down robotic snake.



A negative result:

Theorem. (A. - Yatchak, 2012) If the snake in a grid is not pinned down, the state complex is not always $CAT(0)$.

Open question. Which robots give $CAT(0)$ cube complexes?

Application 5. Moving (some) robots efficiently.

Motivation:

Algorithm. (Owen-Provan 09) A polynomial-time algorithm to find the geodesic between trees T_1 and T_2 in the space of trees \mathbf{T}_n .

($\sqrt{2}$ -approx.: Amenta 07, exp.: GeoMeTree 08, GeodeMaps 09)

This allows us to

- find distances between trees
- “average” trees.

Application 5. Moving (some) robots efficiently.

We use the PIP (“remote control”) of X to get:

Algorithm. (A. - Owen - Sullivant 12, A - Baker - Yatchak 14)

An algorithm to find the geodesic between points p and q in **any** CAT(0) cube complex X .

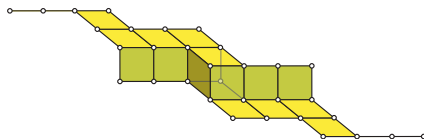
We do this for four metrics:

- Euclidean length
- Time
- Number of moves.
- Number of sets of simultaneous moves.

This allows us to

- navigate the state complex of any reconfiguration system
- find the optimal robot motion between two positions.

(Computer/robotic implementation?)



much as gr acias

Los primeros dos artículos y esta presentación están en:

Advances in Applied Mathematics **48** (2012) 142-163.

SIAM J. Discrete Math. **28-2** (2014), pp. 986-1007

<http://arxiv.org/abs/1101.2428>

<http://arxiv.org/abs/1211.1442>

<http://math.sfsu.edu/federico>